## B.Tech. (Sem. - $7^{\text {th }} / 8^{\text {tr }}$ )

# FORMAL LANGUAGE \& AUTOMATA THEORY <br> SUBJECT CODE : CS - 404 <br> Paper ID : [A0481] 

[Note : Please fill subject code and paper ID on OMR]

## Time : 03 Hours

Maximum Marks : 60

## Instruction to Candidates:

1) Section - A is Compulsory.
2) Attempt any Four questions from Section - B.
3) Attempt any Two questions from Section - C.

## Section - A

Q1)
a) What is context-free grammar?
b) What is Griebach normal form?
c) What do you understand by type-1 grammar?
d) What are palindromes?
e) What do you understand by acceptability of language by turing machine?
f) What is type-2 grammar?
g) Define $\mathrm{LR}(\mathrm{k})$ grammar.
h) Why do natural languages are not formal languages?
i) Represent the set of regular expression of all strings over $(a, b)$ beginning and ending with a.
j) What do you understand by term union of sets?

## Section-B

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(4 \times 5=20)
$$

Q2) Prove $\left(1+00^{*} 1\right)+\left(1+00^{*} 1\right)\left(0+10^{*} 1\right)^{*}\left(0+10^{*} 1\right)=0^{*} 1\left(0+10^{*} 1\right)^{*}$.

Q3) Define turing machine in details.
Q4) Define pushdown automata completely:
Q5) Design a turing machine over $\{1, \mathrm{~b}\}$ which can compute concatenation function over $\Sigma=\{1\}$. If a pair of words (w1,w2) is an input, the output has to be w1w2.

Q6) Prove that grammar $S->0 \mathrm{~A} 2, \mathrm{~A}->1 \mathrm{~A} 1, \mathrm{~A}->1$ is not $\mathrm{LR}(0)$.

## Section - C

$(2 \times 10=20)$
Q7) Write note on universal turing machine and modification of basic model of turing machine.

Q8) Construct a pda accepting empty store of languages in each case :
(a) $\left\{a^{n} b^{m} a^{n} \mid m, n>=1\right\}$.
(b) $\left\{a^{n} b^{2 n} \mid n>=1\right\}$.
(c) $\left\{\mathrm{a}^{m} \mathrm{~b}^{m} \mathrm{c}^{n} \mid \mathrm{m}, \mathrm{n}>=1\right\}$.
(d) $\left\{\mathrm{a}^{m} \mathrm{~b}^{n} \mid \mathrm{m}>\mathrm{n}>=1\right\}$.

Q9) Construct a regular grammar $G$ generating the regular set represented by $P=a^{*} b(a+b)^{*}$.

