

THEORY OF COMPUTATION

Q1. Choose the correct statement.

The set of all strings over an alphabet $S = \{0,1\}$ with the concatenation operator for strings

- a) does not form a group
- b) forms a noncommutative group
- c) does not have a right identity
- d) forms a group if the empty string is removed from S^*

Q2. Consider the set of all strings S^* over an alphabet $S = \{a,b\}$ with the concatenation operator for strings, and

- a) the set does forms semigroup
- b) the set does not form a group
- c) the set has a left and right identity
- d) the set forms a monoid

Q3. Consider the set of all strings S^* over the alphabet $S = \{a,b,c,d,e\}$ with the concatenation operator for strings.

- a. the set has a right identity and forms a semigroup
- b. the set has a left identity and forms a monoid
- c. the set does not form a commutative group but has an identity
- d. the set does not form a semigroup with identity

Q4. Nobody knows yet if $P = NP$. Consider the language L defined as follows:

$L = (0+1)^*$ if $P = NP$

And

$L = j$ otherwise

Which of the following statements is true?

- a) L is recursive
- b) L is recursively enumerable but not recursive
- c) L is not recursively enumerable
- d) Whether L is recursive or not will be known after we find out if $P = NP$

Q5. Consider the language defined as follows

$L = \{a^n b^n \mid n \geq 1\}$ if $P = NP$

And

$L = \{ww \mid w \in (a+b)^+\}$ otherwise

Which of the following statements is true?

- a) L is recursive but not context sensitive
- b) L is context sensitive but not context free
- c) L is context sensitive
- d) What L is will be known after we resolve the $P = NP$ question

Q6. Consider the language defined as follows

$L = (0+1)^*$ if the CSLs are closed under complement

And

$L=(0^*1)^*0^*$ if $P=NP$

And

$L=(10^*)1^*$ if P is not the same as NP

Which of the following statements is true?

- a) L is always a regular set
- b) L does not exist
- c) L is recursive but not a regular set
- d) What L is will be known after the two open problems $P = NP$ and the closure of CSLs under complement are resolved

Q7. Consider the language defined as follows

$L=(0+1)^*$ if man goes to Mars by 2020AD

And

$L=0^*$ if man never goes to the Mars

Which of the following is true?

- a. L is context free language but not recursive
- b. L is recursive
- c. Whether L is recursive or not will be known in 2020AD
- d. L is a r.e. set that is not regular

Q8. Given an arbitrary context free grammar G , we define L as below.

$L=(0+1)^*$ if G is ambiguous

And

L=j if G is not ambiguous

a. L is a context-free language

b. L is recursive but not r.e.

c. What L is depends on whether we can determine if G is ambiguous or not

d. What L is is undecidable

Q9. Given an arbitrary turing machine M and a string w we define L as below.

L=(0*1)*0* if M halts on w

And

L=(0*1*)* if M does not halt on w

a. The type of L is undecidable because of the halting problem

b. L is a context-sensitive language

c. L is recursively enumerable and not context-free

d. L is context sensitive and not regular

Q10. Consider the language L defined below

L=(0+1)* if P=NP

And

L=(aⁿbⁿ|n>=1} otherwise

a. Whether L is a regular set that is not context-free will be known after we resolve the P=NP question.

b. Whether L is context-free but not regular will be known after we resolve the P=NP question

c. L is context-sensitive

d. L is not recursive

Q11. It is undecidable if two cfls L1 and L2 are equivalent. Consider two cfls L1 and L2 and a language defined as follows

$L = \{a^n b^n c^n \mid n \geq 234\}$ if $L1 = L2$

And

$L = \{a^n b^n c^n d^n \mid n \geq 678\}$ otherwise

a. L is recursive

b. L is context-free

c. We can never say anything about L as it is undecidable

d. L is regular

Q12. At present it is not known if NP is closed under complementation.

Consider L defined as below

$L = \{w wR w \mid w \text{ in } (0+1+2)^* \text{ and } wR \text{ is the reverse of } w\}$ if NP is closed under complement

And

$L = \{a^n b^n c^n d^n e^n \mid n \geq 34567\}$ otherwise

- a) L is recursive
- b) L is context-free and not context-sensitive
- c) L is recursively enumerable but not recursive
- d) We will be able to say something about L only after we resolve the NP complementation issue

Q14. Nobody knows if P=NP at present. Consider a language L as defined below

L=(0+1)* if satisfiability is in P

L=(0*1)0* if satisfiability is not in P

L=(1*0)1* if 3-sat is in P

L=(0*1*)* if 3-sat is not in P

L=(0*1*0*1*)* if 0/1 knapsack problem is in P

L=(1*0*1*0*)* if 0/1 knapsack problem is not in P

L=(0*(00)*(1*11*)) * if max-clique problem is in P

L=(0*(00)*(1*11*)) * if node-cover problem is not in P

L=(0*1*)*** (010)* if edge-cover problem is not in P**

L=(0* + 1* + (00)* + (11))* (0100101010)* if the chromatic problem is not in P

What can we say about the string 0000111100001111=x

- a) x is always in L
- b) whether x is in L or not will be known after we resolve P=NP
- c) the definition of L is contradictory
- d) x can never be in L

Q15. An arbitrary turing machine M will be given to you and we define a language L as follows

$L=(0+00)^*$ if M accepts at least one string

$L=(0+00+000)^*$ if M accepts at least two strings

$L=(0+00+000+0000)^*$ if M accepts at least three strings

$L=(0+00+000+---+0^n)$ *if M accepts at least n-1 strings

Choose the correct statement.

- a) We cannot say anything about L as the question of whether a turing machine accepts a string is undecidable
- b) L is context-sensitive but not regular
- c) L is context-free but not regular
- d) L is not a finite set

Q16. We are given two context-free languages L1 and L2 and L defined as below

a) $L=(0+1)^*$ if $L1=L2$

b) $L=((0+00+000)^*(1+11+111)^*)^*$ if L1 is contained in L2

c) $L=((0(10)^*)^*(1(01)^*)^*)^*$ if L2 is contained in L1

d) $L=(00+11+0+1)^*(0+00+000)^*$ if L1 and L2 are incomparable

a) As all the conditions relating to L1 and L2 are undecidable we cannot say anything about L

- b) L is recursively enumerable
- c) L is recursive but not context-sensitive
- d) L is context-sensitive but not context-free
- e) L is context-free but not regular

Q17. It is undecidable if an arbitrary cfl is inherently ambiguous. We are given a cfg G and the language L is defined as below

$L = (0+1)^*01(0+1)^* \cup 1^*0^*$ if L(G) is inherently ambiguous

$L = (0+1)^*10(0+1)^* \cup 0^*1^*$ if L(G) is not inherently ambiguous

Choose the incorrect statement

- a) L is regular
- b) L is context-free
- c) L is context-sensitive
- d) The above choices can be resolved only if we know if L(G) is inherently ambiguous or not

Q18. We are given an arbitrary turing machine M and define the language L as below

$L = (0^*+1^*)^*$ if M halts on blank tape

$L = (0+1)^*$ if M ever prints a 1

$L = (0^*+1)^*$ if M ever enters a designated state q

$L = ((0+1+00+11+000+111)+)^*$ if M accepts an infinite set

$L = 0^*(10^*)^*$ if M accepts a finite set

$L = 1^*(01^*)^*$ if M accepts exactly 45 strings

Choose the correct statement with reference to the string $x=00000111111000000111111$

a) x is in L

b) x is not in L

c) we can never decide if x is in L as all the problems of the turing machine are undecidable

d) whether x is in L depends on the particular turing machine M

Q19. We are given a language L defined as follows

$L=(0+1)^*$ if the Hamiltonian circuit problem is in P

$L=(0^*1^*+0)^*$ if the Traveling salesman problem is not in P

$L=(0^*1^*1)^*0^*$ if the bin packing problem is in P

a) the definition of L is contradictory

b) What L is will be known after we resolve the P=NP question

c) L is a finite set

d) The string 01010101001010110010101 is in L

Q20. The intersection of two cfls can simulate a turing machine computation. We are given two cfls L1 and L2 and define the language L as below

a) $L=(00)^*$ if the intersection of L1 and L2 is empty

b) $L=((0(00)^*)(0(00)^*))^*$ if L1 is regular

c) $L=(00+0000+000000)^*$ if L2 is not regular

d) $L=(00)^*+(0000)^*$ if the complement of L1 is a cfl

a) L is a finite set

b) L is a regular set

c) L is undecidable

d) L is recursive but not context-free